



THE UNIVERSITY
OF BRITISH COLUMBIA

Approximate Normalization for Gradual Dependent Types

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ICFP 2019

Our Contributions

GDTL:

Gradual Dependently Typed Language

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- Full spectrum, universe hierarchy

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GDTL:

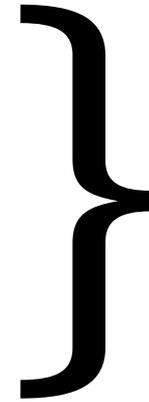
Gradual Dependently Typed Language

- Full spectrum, universe hierarchy
- Can replace any type or term with ?
- Embeds fully typed & untyped calculi
- Decidable typechecking

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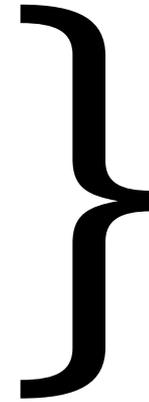


This talk

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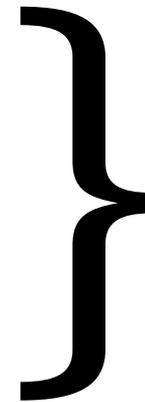
This talk

- Proof of gradual type safety

GDTL:

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- Proof of gradual type safety
- Gradual Guarantees (Siek et al 2015):
reducing precision of term won't
create new static or dynamic
failures

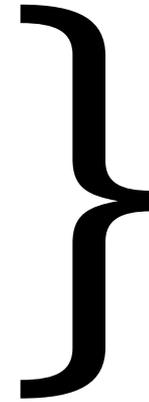


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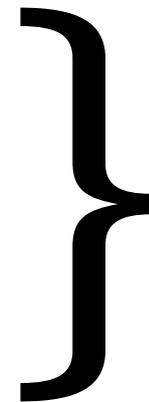
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This talk



See paper

Why Gradual Dependent Types?

Dependent Types: Pain and Promise

Expectation

```
$> compile ./myprogram  
>> 0 bugs detected!
```

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Reality

```
$> compile ./myprogram  
>> Type mismatch between  
      Vec Nat (m+n)  
      and  
      Vec Nat (n+m)
```

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- Popular for proof assistants

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Reality

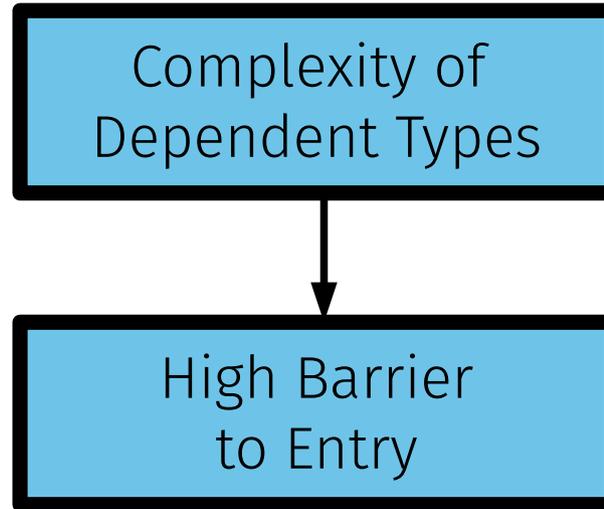
```
$> compile ./myprogram  
>> Type mismatch between  
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      and  
      Vec Nat (n+m)
```

- Popular for proof assistants
- Not popular for programming

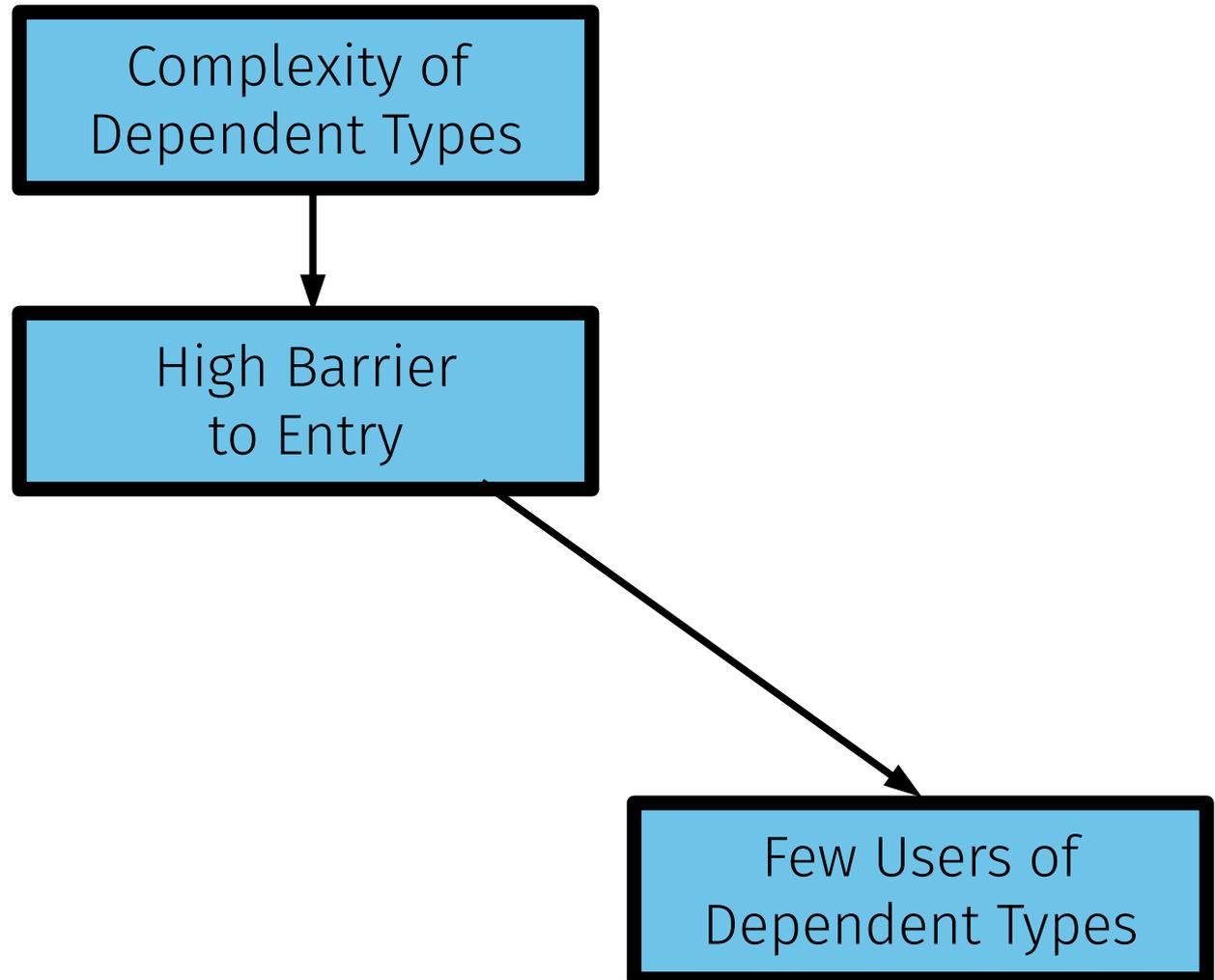
A Vicious Cycle

Complexity of
Dependent Types

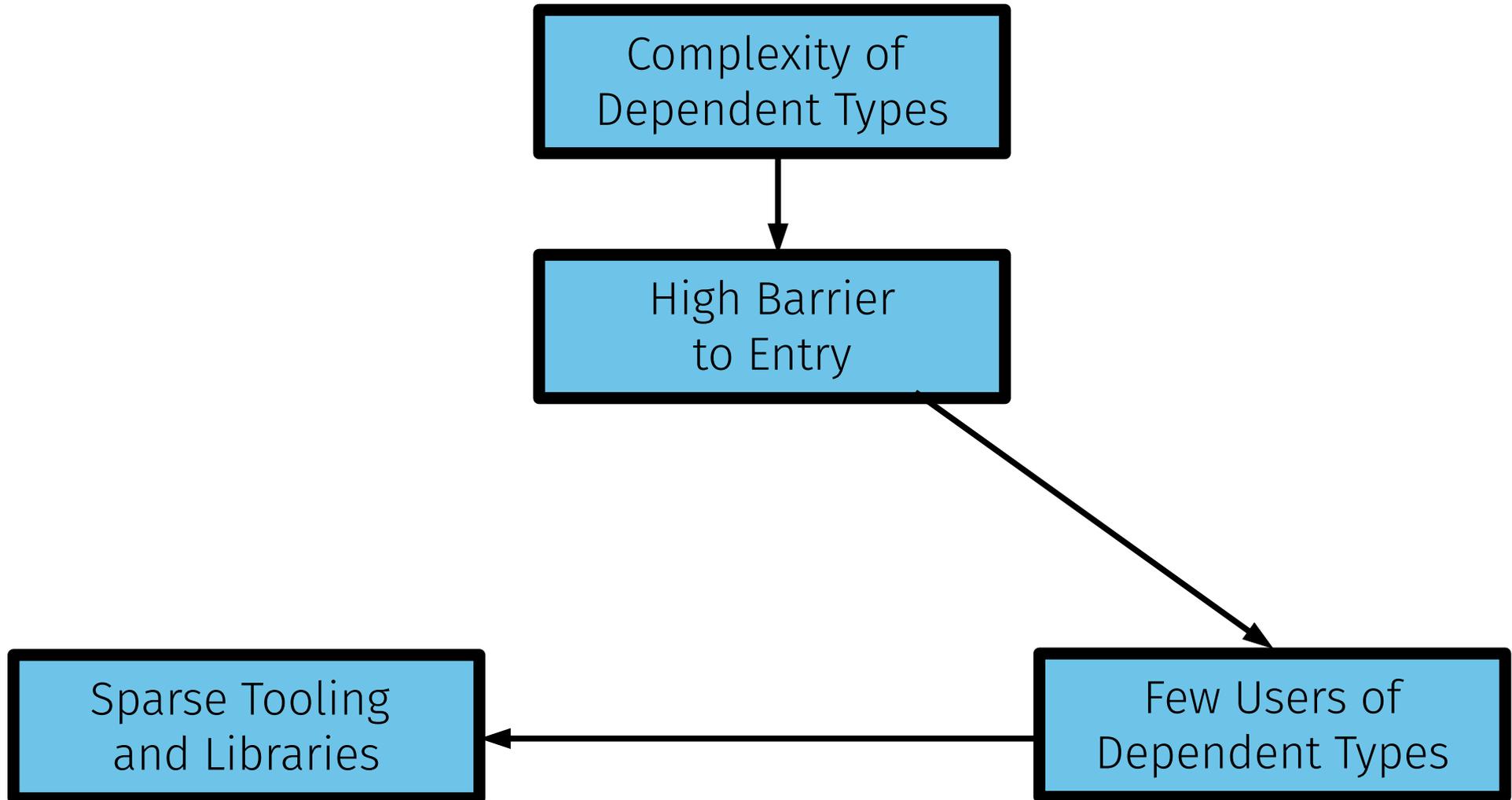
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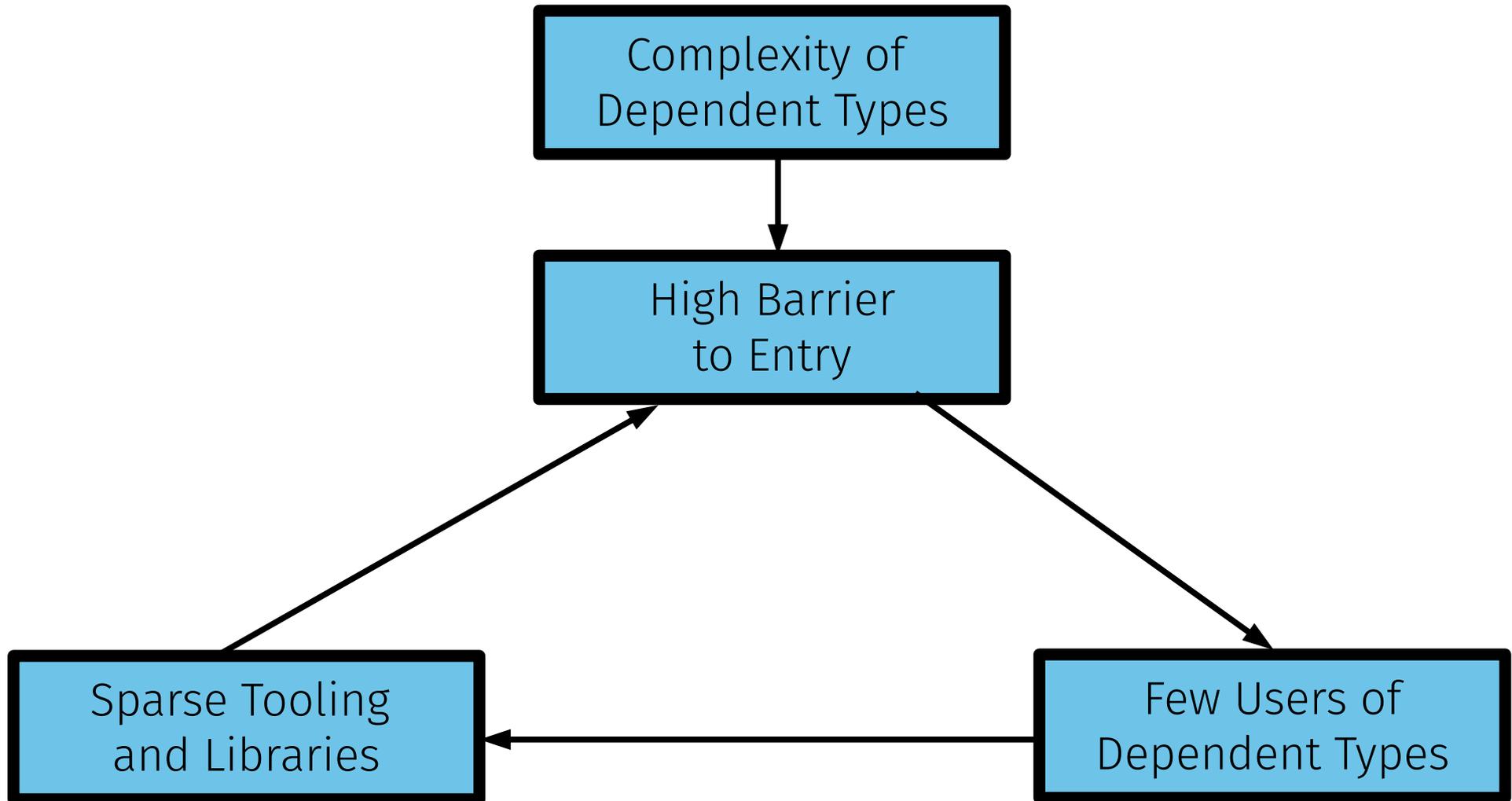
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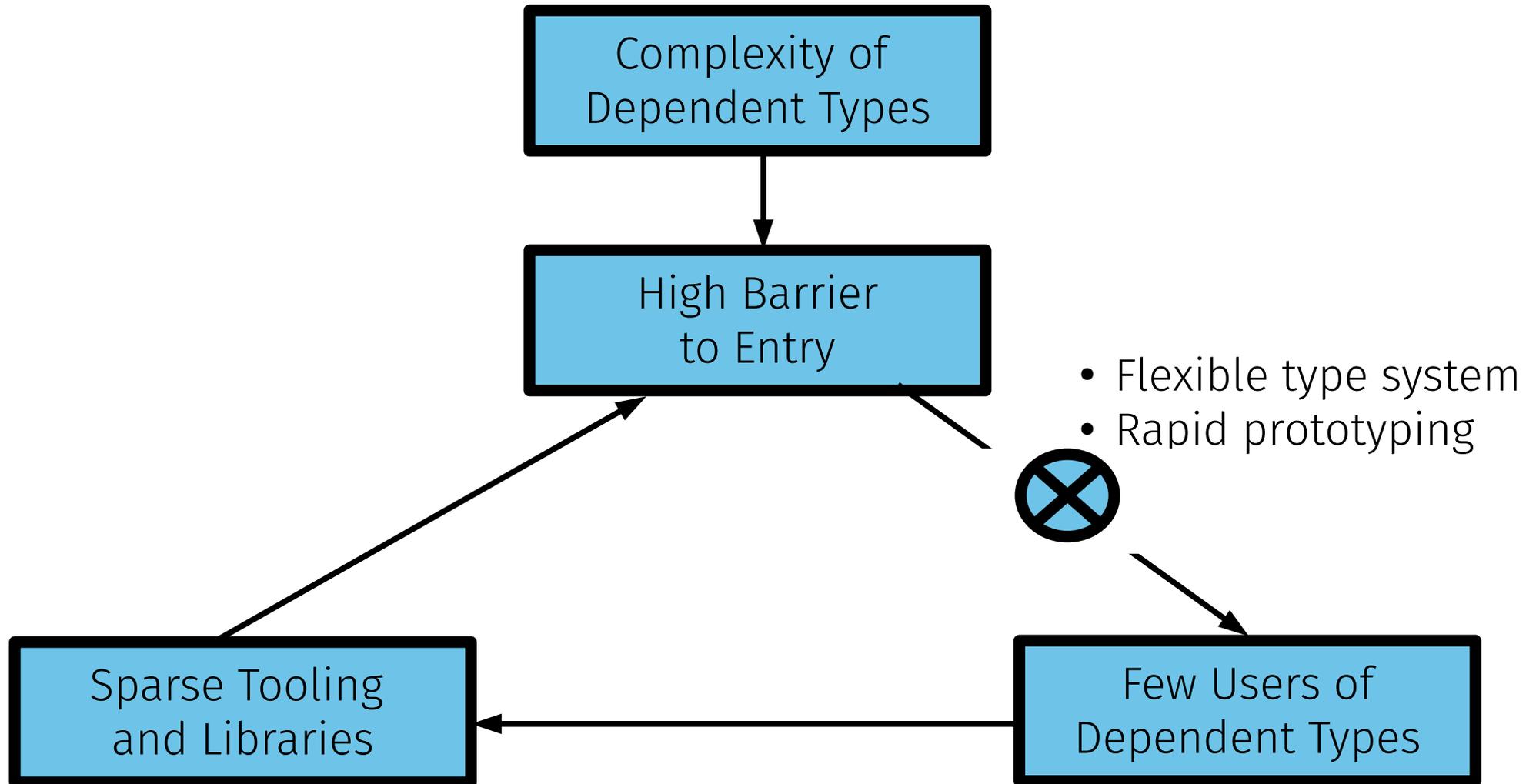
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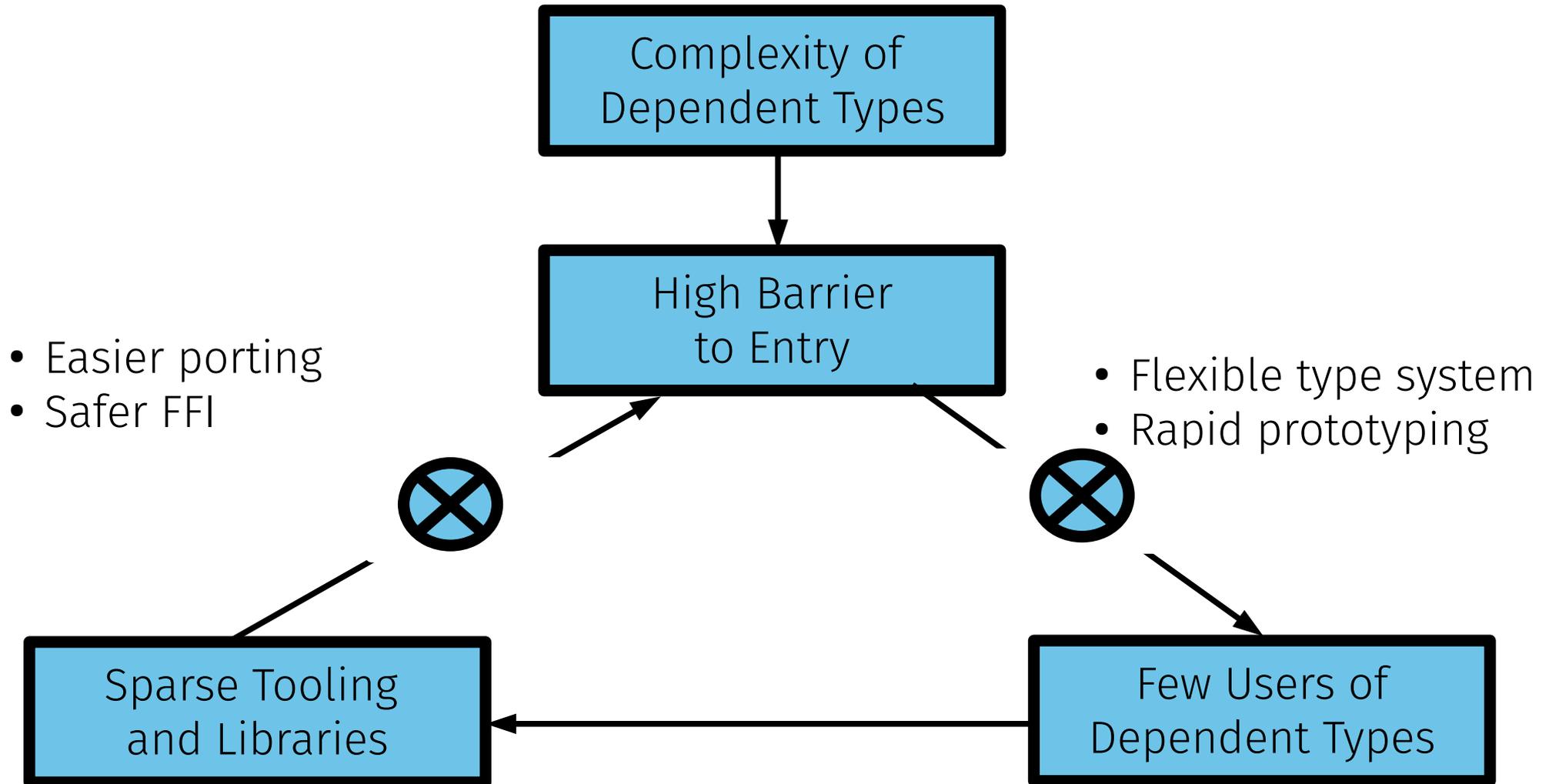
A Vicious Cycle



A Vicious Cycle



A Vicious Cycle



Goals For Gradual Dependent Types

Motivation: Lists vs. Vectors

Motivation: Lists vs. Vectors

```
data List a
  where
    Nil : List a
    Cons : a
          -> List a
          -> List a

head : List a -> a
```

Motivation: Lists vs. Vectors

No size knowledge
in type

```
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Error on Nil

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```
data List a
  where
    Nil : List a
    Cons : a
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head : List a -> a
```

```
data Vec (a : Type) (n : Nat)
  where
    Nil : Vec a 0
    Cons : a
           -> Vec a n
           -> Vec a (n + 1)

head : Vec a (n + 1) -> a
```

Error on Nil

Motivation: Lists vs. Vectors

No size knowledge
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data List a
  where
    Nil : List a
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head : List a -> a
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Error on Nil

Length in *type index*

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data Vec (a : Type) (n : Nat)
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    Nil : Vec a 0
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head : Vec a (n + 1) -> a
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Motivation: Lists vs. Vectors

No size knowledge
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data List a
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  where
    Nil : Vec a 0
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head : Vec a (n + 1) -> a
```

Won't typecheck for Nil

Porting Code to Dependent Types

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```
sort : List Int -> List Int
```

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sort Nil = Nil
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```
sort (Cons head tail) =
```

Porting Code to Dependent Types

```
sort : List Int -> List Int  
  
sort Nil = Nil  
  
sort (Cons head tail) =  
  sort (filter (<= head) tail))
```

Porting Code to Dependent Types

```
sort : List Int -> List Int  
  
sort Nil = Nil  
  
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  ++ [head]
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filter : (Int -> Bool)  
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```

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      Vec Int n
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      Vec Int n      Vec Int n
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      Vec Int n      Vec Int n  
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sort Nil = Nil ✓
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Vec Int n    Vec Int n  
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Need proof
that
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Vec Int n    Vec Int n  
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filter : (Int -> Bool)  
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Porting Code to Dependent Types

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          Vec Int n    Vec Int ___
```

How To Solve?

How To Solve?

Static Dependent Types		

How To Solve?

Static Dependent Types	Existential Types, Inductive Proof	

The Static Approach

The Static Approach

`rewriteFilterLength :`

The Static Approach

```
rewriteFilterLength :  
  (v : Vec Int n)
```

The Static Approach

```
rewriteFilterLength :  
  (v : Vec Int n)  
  -> (p : Int -> Bool)
```

The Static Approach

```
rewriteFilterLength :  
  (v : Vec Int n)  
-> (p : Int -> Bool)  
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    (length (filter p v)  
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```

Relies on induction, commutativity, etc.

Gradual Proof Terms

```
sort : Vec Int n -> Vec Int n
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filter : (Int -> Bool)
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Static Dependent Types	Existential Types, Inductive Proof	x Significant effort required

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Static Dependent Types	Existential Types, Inductive Proof	X Significant effort required
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How To Solve?

Static Dependent Types	Existential Types, Inductive Proof	 Significant effort required
Non-dependent Gradual Types	<code>filter</code> returns ? unknown <u>type</u>	 Can have non-list return
Gradual Dependent Types	<code>filter</code> returns <code>Vec Int</code> ? unknown <u>length</u>	 Precise in type, flexible in length!

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Static Dependent Types	Existential Types, Inductive Proof	✗ Significant effort required
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Gradual Dependent Types	<code>filter</code> returns <code>Vec Int ?</code> unknown <u>length</u>	✓ Precise in type, flexible in length!

Our approach!

The GDTL Solution

```
sort : Vec Int n -> Vec Int n
```

```
sort Nil = Nil
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```



?+1+? evals
to ?, is
consistent
with n

```
filter : (Int -> Bool)  
        -> Vec Int n -> Vec Int ?
```

Gradual Type Safety

Gradual Type Safety

Gradual Type Safety

head : Vec a (n+1) -> a

$\text{head} : \text{Vec } a \ (n+1) \rightarrow a$

$x : \text{Vec } a \ 0$

$\text{head} : \text{Vec } a \ (n+1) \rightarrow a$

$x : \text{Vec } a \ 0$

$\text{theHead} = \text{head } x$

Gradual Type Safety

`head : Vec a (n+1) -> a`

`x : Vec a 0`

`theHead = head x`

- Does not typecheck

$\text{head} : \text{Vec } a \ (n+1) \rightarrow a$

$x : \text{Vec } a \ 0$

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$\text{head} : \text{Vec } a \ (n+1) \rightarrow a$

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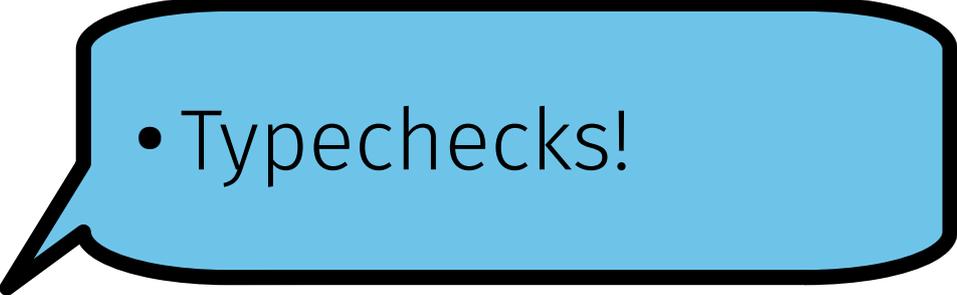
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Gradual Type Safety

`head : Vec a (n+1) -> a`

`x : Vec a ?`

`theHead = head x`

- 
- Typechecks!

$\text{head} : \text{Vec } a \ (n+1) \rightarrow a$

$x : \text{Vec } a \ ?$

$\text{theHead} = \text{head } x$

$x \mapsto \text{Nil}$

Gradual Type Safety

$\text{head} : \text{Vec } a \ (n+1) \rightarrow a$

$x : \text{Vec } a \ ?$

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$x \mapsto \text{Nil}$

- Runtime error

Gradual Type Safety

`head : Vec a (n+1) -> a`

`x : Vec a ?`

`theHead = head x`

`x ↦ Nil`

`x ↦ Cons 1 Nil`

Gradual Type Safety

`head : Vec a (n+1) -> a`

`x : Vec a ?`

`theHead = head x`

`x` \mapsto

- Runs successfully

`x` \mapsto `Cons 1 Nil`

Filling in the Proof

```
sort : Vec Int n -> Vec Int n
```

```
sort Nil = Nil
```

```
sort (Cons head tail) =
```

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  sort (filter (<= head) tail))  
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Filling in the Proof

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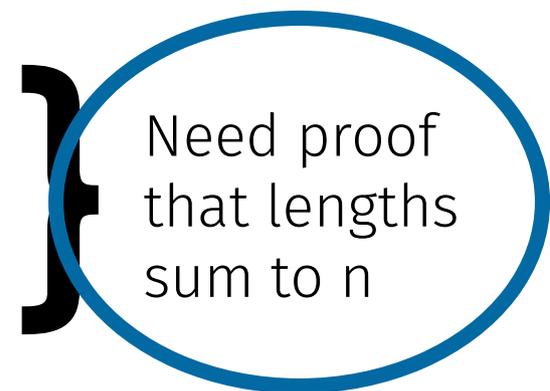
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```
filter : (Int -> Bool) -> Vec Int n -> Vec Int ?
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Gradual Proof Terms

```
rewriteFilterLength :  
  (v : Vec Int n)  
-> (p : Int -> Bool)  
-> Vec Int  
    (length (filter p v)  
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-> Vec Int n
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Gradual Proof Terms

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rewriteFilterLength = ?

Gradual Proof Terms

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```

rewriteFilterLength = ? } Like
 } Idris/Agda
 } typed holes

Gradual Proof Terms

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```

```
filter : (Int -> Bool)  
        -> Vec Int n -> Vec Int ?
```

This code typechecks *and runs!*

Semantics of ? in GDTL

- ? has type ?, can use at any type

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- Eliminating ? produces ?

- ? has type ?, can use at any type
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- Runtime checks ensure safety

Semantics of ? in GDTL

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subst : $a = b \rightarrow P a \rightarrow P b$

Semantics of ? in GDTL

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`subst` : $a = b \rightarrow P a \rightarrow P b$

`badProof` : $0 = 1$

`badProof` = ?

Semantics of ? in GDTL

- ? has type ? , can use at any type
- Eliminating ? produces ?
- Runtime checks ensure safety

`subst : a = b -> P a -> P b`

`badProof : 0 = 1`

`badProof = ?`

`head ((subst badProof nil) :: Vec Int 1)`

Semantics of ? in GDTL

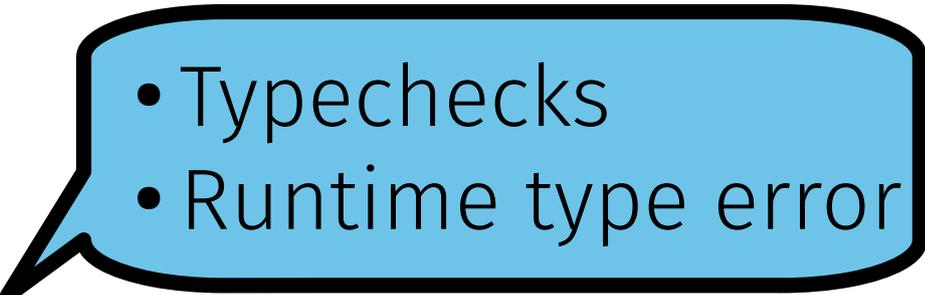
- ? has type ?, can use at any type
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`subst : a = b -> P a -> P b`

`badProof : 0 = 1`

`badProof = ?`

`head ((subst badProof nil) :: Vec Int 1)`

- 
- Typechecks
 - Runtime type error

GDTL: A Gradual Dependently Typed Language

Gradual Dependent Types

Gradual Dependent Types

Statics + Dynamics mostly using
Abstracting Gradual Typing (Garcia et. al. 2016)

Gradual Dependent Types

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Main extensions:

Gradual Dependent Types

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Main extensions:

Type/Term Overlap

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Main extensions:

Type/Term Overlap \longrightarrow ? as unknown
type *and* term

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Type Indices

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Main extensions:

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Type Indices \longrightarrow ? as type index

Gradual Dependent Types

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Main extensions:

Type/Term Overlap \longrightarrow ? as unknown
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Type Indices \longrightarrow ? as type index

Proof term

Gradual Dependent Types

Statics + Dynamics mostly using
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Main extensions:

Type/Term Overlap \longrightarrow ? as unknown
type *and* term

Type Indices \longrightarrow ? as type index

Proof term \longrightarrow ? as a term
at runtime

What's the Catch?

Dependent Types

Dependent Types

Evaluate terms
at compile time

Dependent Types

Evaluate terms
at compile time

Strongly
normalizing

Dependent Types

Evaluate terms
at compile time

Strongly
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Failure free

Dependent Types

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Gradual Types

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Gradual Types

Evaluating has
effects

Dependent Types

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Gradual Types

Evaluating has
effects

Can diverge
i.e. $\lambda(x : ?). x x$

Dependent Types

Evaluate terms
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Gradual Types

Evaluating has
effects

Can diverge
i.e. $\lambda(x : ?). x x$

Type errors in
evaluation

Key Idea: Approximate Normalization

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Exploit the phase distinction:

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Exploit the phase distinction:

Compile-time Normalization		

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Compile-time Normalization	Always terminates	

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Key Idea: Approximate Normalization

Exploit the phase distinction:

Compile-time Normalization	Always terminates	Approximate results
Runtime Evaluation	May diverge	Exact results

Compile-Time - Approximation #1: Termination

Based on *Hereditary Substitution*
(Watkins et al 2003)

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Types structurally decreasing?

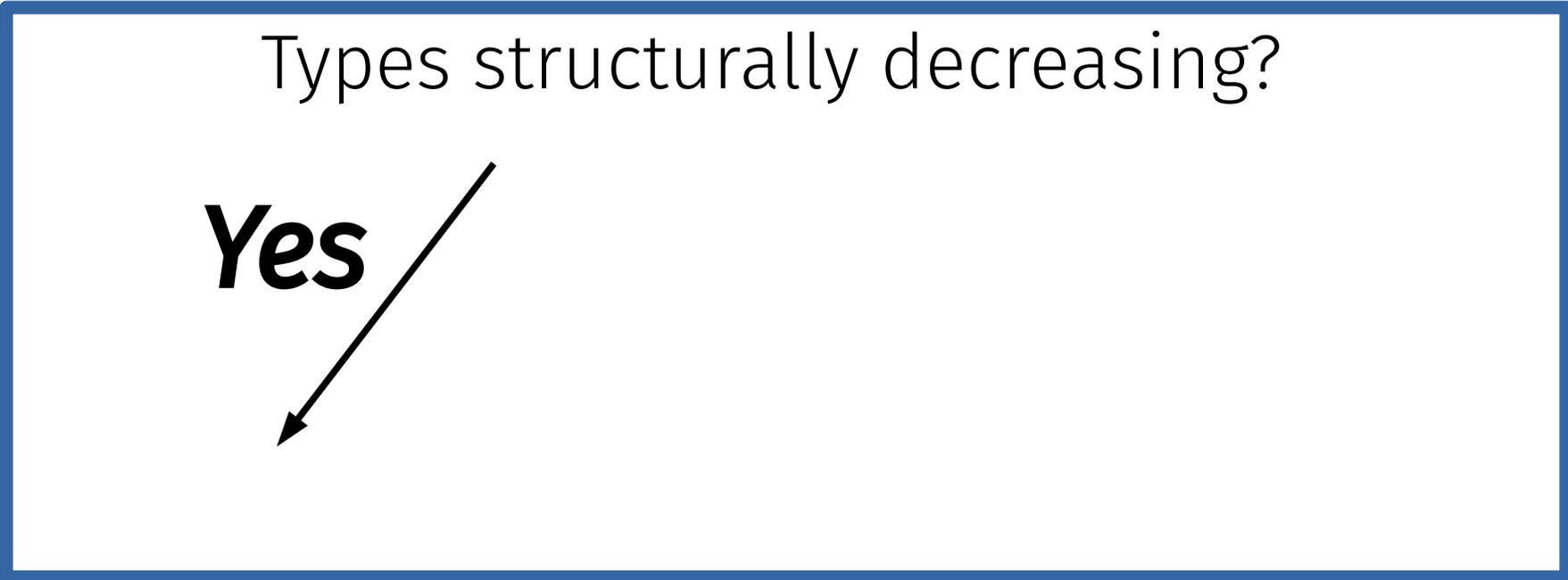
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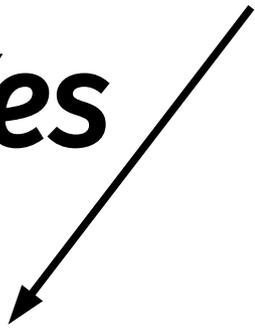
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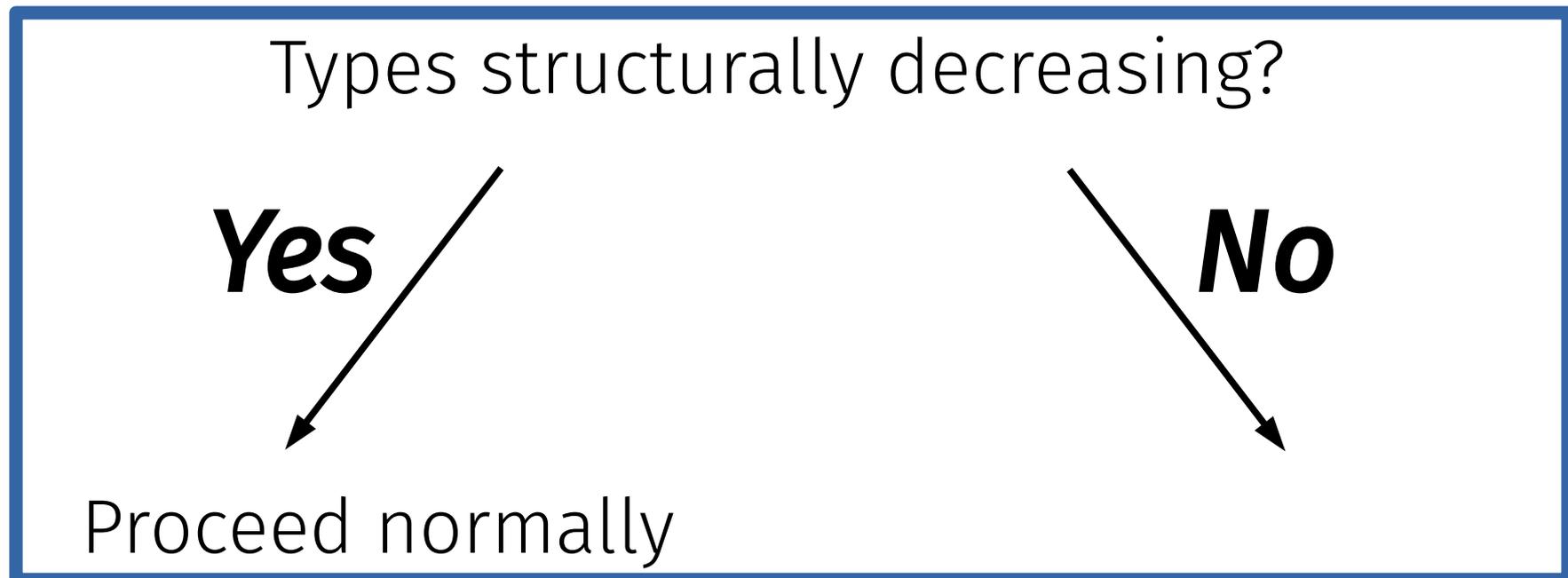
Proceed normally

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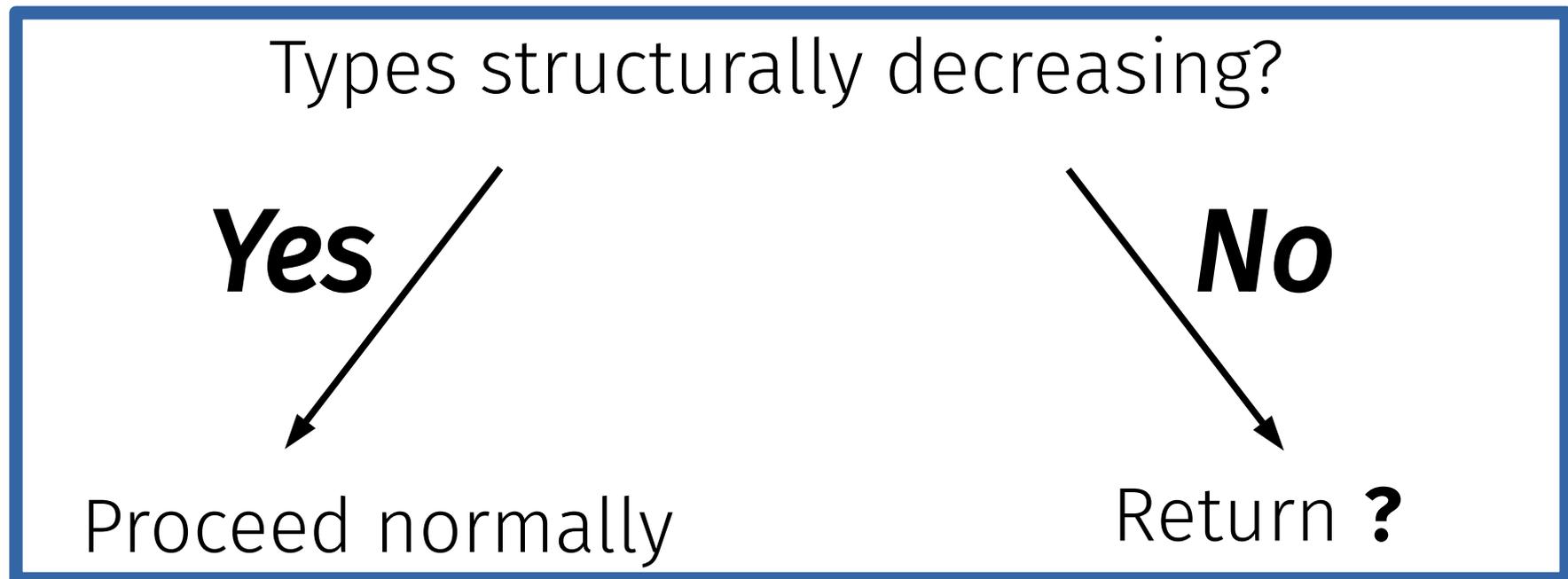
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$$(\lambda x. x(\lambda y. yy))(\lambda z. zz)$$

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 - Runtime error if meet does not exist

Wrapping Up

What We Built

GDTL:

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Gradual Dependently Typed Language

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Future Work

- Inductives and Pattern Matching

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- Inductives and Pattern Matching
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- Blame and Error Reporting
- Eventual Goal: Idris frontend

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